

The design and implementation of embedded Domain-Specific Languages

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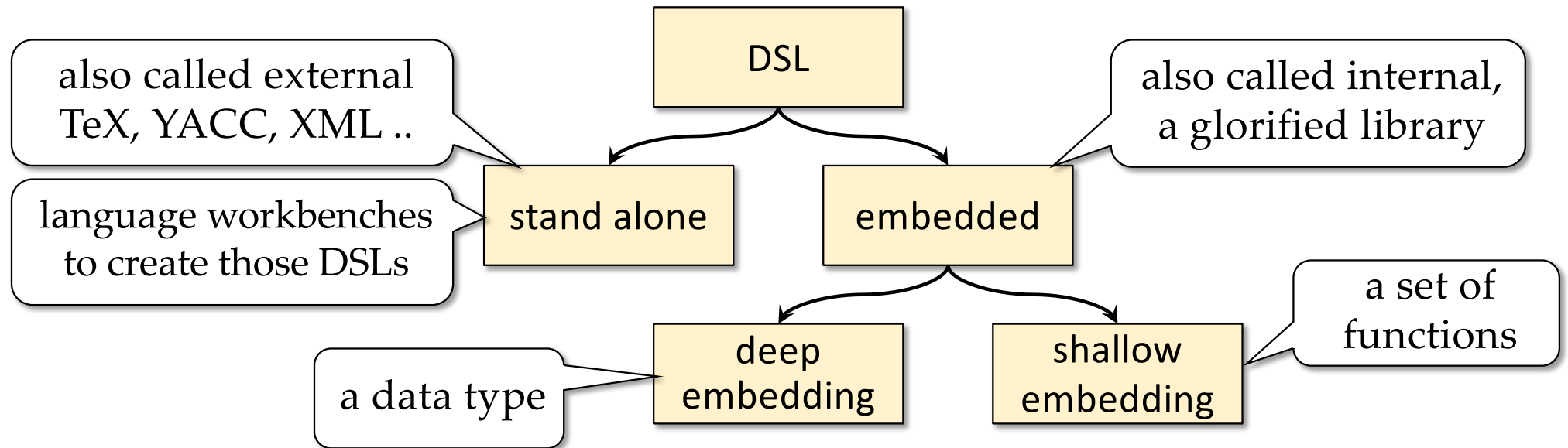


Domain-Specific Languages, DSLs

- DSL is specialized programming language for a particular application domain
 - e.g. **HTML** for webpages, **TeX** for type setting, **SQL** for database queries, **SVG** for graphics, GraphViz's **dot** for visualisation of graphs, **YACC** parser construction, **QuickCheck** for properties in model-based testing, ...
- in contrast to a general-purpose language (GPL)
 - e.g. Haskell, Java, Erlang, Scala, C/C++, ...
- Why a DSL?
 - **make programming and maintenance easier and cheaper**
 - DSL offers abstraction level tailored to problem domain
 - DSL implementation avoids repeated work



implementation strategies for DSLs



- we focus on embedded DSLs
to reuse host language implementation, libraries and tools
- Two concepts: **views**, **DSL constructs**



implementation goals for eDSL

1. Multiple views

excludes shallow embedding (functions)

- execution (simulation), and pretty printing, code generation, ...
- adding views should not break anything in existing views and DSL programs

2. Type safety

- eDSL has same level of trust as host language
- no runtime type errors
- eDSL constructs can be overloaded (like equality, comparison, addition)

3. Safe identifiers

- No runtime errors (so no strings or numbers as identifiers)

4. Extendable DSL

- extend DSL without breaking existing code (programs and views)

Wadler's expression problem



eDSL and running example



- eDSL to control peripherals (temperature sensor, heating device, ...)
- eDSL is functional and task-oriented

program1 =

define sens

define heat

Loop (

Read sens >>=. \temp ->

Post (temp < 19) heat

)

how to define
these identifiers?

host language
our eDSL

- simulation view in the browser of this program



Temperature:	<input type="text" value="7"/>	✓
Heating:	True	

Temperature:	<input type="text" value="19"/>	✓
Heating:	False	



1. multiple views
2. type safety
3. safe identifiers
4. extendable DSL

Deep embedding



our example eDSL in a deep embedding

- Deep embedding
 - language constructs as constructors *of* an ADT
 - views (interpretations) as functions *over* the ADT
 - hosted in Clean, a language similar to Haskell.
- views of the eDSL we will make
 - pretty printing (code generation is essentially equal)
 - simulation in web-browser (using the iTask framework)

Nicholas Rinaudo's
talk of day 1

Plasmeijer's
tutorial of day 1



Expression DSL

```
:: Expr a
   = Lit a
   | E.b: Less (Expr b) (Expr b) & <, type b -> Expr Bool
   ...
```

argument fixes
type of expression

class restrictions
on type variable

expression yields
a Boolean: GADT

// add any operation you need in the DSL

- we need quite some machinery to make this safe and looking good
 - GADT, existential quantified type variables, class restrictions in types, infix constructors
 - explicit projection pairs/bimaps can replace GADTs
 - languages like Haskell and Clean are made for this



Simple views on the expression DSL

```
print :: (Expr a) -> String | type a
print (Lit a)      = toString a
print (Less l r) = print l +++ "<" +++ print r
print ...          = ...
```

```
eval :: (Expr a) -> a
eval (Lit a)      = a
eval (Less l r) = eval l < eval r
eval ...          = ...
```

using GADTs



1. multiple views
2. type safety
3. safe identifiers
4. extendable DSL

Higher-order abstract syntax, HOAS



Higher-order abstract syntax (HOAS)

- Use functions in the host language to define variables
 - Nameless functions (lambdas) are useful but not required
- use higher-order programming
 - What would you do with a value if you had it as some point (monads without do)
`do { x <- compute; f x; }`
this is a shorthand notation for
`compute >>= \x -> f x`
- We can still print, evaluate, simulate, ...
but transformations are more tricky



expressions in our task-oriented eDSL

```
w2Deep :: Work Bool
```

```
w2Deep =
```

```
  DefSens "Temperature" 7 \sens ->
```

```
  DefActr "Heating" False \heat ->
```

```
  Loop (
```

```
    Read sens >>=. \temp ->
```

```
    Post ( Less temp (Lit 19) ) heat
```

```
  )
```

host language

our eDSL



our eDSL in a deep embedding

`:: Sens a = ...`

`:: Actr a = ...`

`:: Work a`

`= E.b: DefSens String b ((Sens b) -> Work a) & type b`

`| E.b: DefActr String b ((Actr b) -> Work a) & type b`

`| Read (Sens a)`

`| Post (Expr a) (Actr a)`

`| E.b: (>>=.) infixl 1 (Work b) ((Expr b)->Work a) & type b`

`| Loop (Work a)`

`| (.|||.) infixr 3 (Work a) (Work a)`

quantified
type variable

HOAS for safe
identifiers

class
restriction

infix constructor

sequencing work

repeating work

parallel work



Deep embedding: simulation

host language
our eDSL
the iTask eDSL

```
sim :: (Work a) -> Task a | type a // big step
sim w = case w of
  w >>=. f = sim w >>- \a -> sim (f (Lit a ))
  Loop w    = sim (w >>=. \_ -> Loop w)
  v .||. w = sim v -||- sim w
  DefSens n v f =
    withShared v \sds ->
      (Label n @>> updateSharedInformation [] sds) ||-
      sim (f (Sens (Sds sds)))
  Read (Sens (Sds s)) = get s
  ...
```



Deep embedding: simulation

```
sim :: (Work a) -> Task a | type a // big step
```

```
sim w = case w of
```

Temperature:

7



Heating:

True

running in browser

```
sim (v >>= . \a -> eval c) u w)
```

Temperature:

19



Heating:

False

```
withShared v \sds ->
```

```
(Label n @>> updateSharedInf
```

```
sim (f (Sens (Sds sds)))
```

```
Read (Sens (Sds s)) = get s
```

Temperature:

18



Heating:

True

...



review of deep embedding

1) multiple views

- just add a function over our datatypes
- clear distinction between host language and eDSL

2) type safety

- type-class system of host language takes care of this
- we use GADTs, extensional quantified type-variables, ..
- intensional analysis is easy

3) safe identifiers

- use higher-order functions (HOAS)

4) extendable DSL 🤔

- change the datatypes, limited compiler support to adjust all views



1. multiple views
2. type safety
3. safe identifiers
4. extendable DSL

class-based embedding, a set of type-constructor classes
tagless-final embedding



expressions

```
class expr v where
```

```
  lit :: a -> v a | type a
```

```
  (<.) infix 4 :: (v b) (v b) -> v Bool | <, type b
```

types are a bit simpler than
in the deep embedding

```
:: Print a = P String      // print view
```

```
instance expr Print where
```

```
  lit a = P (toString a)
```

```
  (<.) a b = P (uP a +++ " <." +++ uP b)
```

actual implementation has
source of fresh identifiers

```
:: Eval a = Eval a        // simulation view
```

```
instance expr Eval where
```

```
  lit a = pure a
```

```
  (<.) a b = (<) <$> a <*> b
```

the *Eval* monad makes this
more concise



our running example

```
:: WorkC v a ::= v (Task a)

w2Class :: WorkC v Bool | dsl v
w2Class =
  defSensC "Temperature" 7 \sens ->
  defActrC "Cold_alarm" False \alarm ->
  loopC (
    readC sens >>=.. \temp ->
    postC (temp <. lit 15) alarm
  )
```

functions instead
of datatypes

```
defSensC Temperature 7 ->
defActrC Cold_alarm False ->
loopC (
  Read Temperature >>=.. \v0 ->
  Post (v0 <. 15 ) Cold_alarm
)
```

Temperature: ✓
Cold_alarm: True

Temperature: ✓
Cold_alarm: False



looping and defining sensors + printing

```
class loopC v :: (WorkC v a) -> WorkC v a | type a
class sens v where
  readC :: (SensC v a) -> WorkC v a | type a
  defSensC :: String b ((SensC v b) -> WorkC v a) -> WorkC v a | ...

instance loopC Print where
  loopC w = P ("loopC " +++ uP w)

instance sens Print where
  readC (SensC n) = P ("Read " +++ uP n)
  defSensC name val f
    = P (print ["defSensC ",name," ",toString val," ->"] +++
        uP (f (SensC (P name))))
```



looping and defining sensors + simulation

```
class loopC v :: (WorkC v a) -> WorkC v a | type a
class sens v where
  readC :: (SensC v a) -> WorkC v a | type a
  defSensC :: String b ((SensC v b) -> WorkC v a) -> WorkC v a | ...

instance loopC Eval where loopC w = w >>=.. \_ -> loopC w

instance sens Eval where
  readC (SensC n) = n >>= pure o get
  defSensC name val f
  = pure (withShared val \sds -> (Label name @>>
    updateSharedInformation [] sds) ||- uE (f (SensC (pure sds)))))
```



review of class-based embedding

1) multiple views

- adding a view is just making a new class instance

2) type safety

- type-class system of host language takes care of this
- types are easier than for deep embedding (datatypes)
- intensional analysis is more tricky

3) safe identifiers

- reuse our technique with higher-order functions

4) extendable DSL

- just add a class to extend the eDSL
- good compiler support to check that required instances exist



overview

	shallow embedding	deep embedding	class-based embedding
	functions	datatypes	classes
evaluation	✓	✓	✓
printing/code	✗	✓	✓
multiple views	✗	✓	✓
eDSL optimization	✗	✓	😓 rather tricky
simple types	✓	🤔 GADTs etc.	✓
good type errors	✓	✓	😓 in terms of classes
extend eDSL	✓	✗ ¹⁾	✓
easy to add views	✗	✓	✓

1) Deep Embedding with Class [doi: 10.1007/978-3-031-21314-4_3](https://doi.org/10.1007/978-3-031-21314-4_3)



what should be part of our DSL

eDSL design



adding constructs to eDSL

Loop w = w >>=. _ -> Loop w

w **Till** f = ...

Read sens **Till** \temp ->
 (Less temp (Lit 19), Post (Lit True) heat)

w **>>|.** f = w >>=. _ -> f

DefTemp f = DefSens "Temperature" 7 f

- YES
 - useful addition
 - other views less simple
- YES
 - useful addition
 - shorter and clearer
- yes ?
 - users might expect it
 - limited added power
 - as macro?
- no
 - limited added value
 - user can define it



adding constructs to eDSL

Loop $w = w \gg= . \setminus _ - \rightarrow \text{Loop } w$

w **Till** f
Read sens
(Less

w **>>|.f** =

DefTemp $f =$ "temperature" 7 f

conditions:

- useful addition to the abstraction to offer
 - adds new power to the language
 - make eDSL programs clearer
 - make eDSL programs shorter
- things people might expect
- use host language as macro language

• YES

- useful addition
- other views less simple

YES

- useful addition
- shorter and clearer

es ?

- users might expect it
- limited added power
- as macro?

• no

- limited added value
- user can define it



Domain Specific Languages - conclusions

- DSLs help you to create and maintain code in specific domain
 - tailor-made abstraction layer, separation of concerns
- embedded (internal) DSLs integrate with host language and reuse tooling
 - external DSLs require parser, type checker, code generator, libraries, tooling, ...
- DSL implementation strategies – with their own advantages and challenges
 - shallow embedding DSL is set of functions
 - deep embedding DSL is datastructure (GADTs, HOAS, ..)
 - class-based embedding DSL is set of type constructor classes
- DSL design
 - do not hesitate to add constructs that improve the creation or maintenance of code
 - keep the DSL as small as possible
 - use the host-language as macro mechanism
- there is much more on eDSLs than we can tell in this tutorial

