



# Embedding Generic Monadic Transformer into Scala.

{Can we merge concurrent programming into mainstream? }

Ruslan Shevchenko, Kyiv, Ukraine.

[<ruslan@shevchenko.kiev.ua>](mailto:ruslan@shevchenko.kiev.ua)

Work:

<https://github.com/rssh/dotty-cps-async>

Industry: [proofspace.id](http://proofspace.id)

Academia: Institute of Software Systems NASU

# Asynchronous control flow - Scala industrial usage.

```
for{
    r1 <- cacheService.get(id) match
        case Some(v) => IO.success(v)
        case None => talkToServer(id).map{v =>
            cacheService.update(id,v)
            v
        }
    r2 <- talkToServer(r1.data)
    result <- if (r2.isOk) then {
        writeToFile(resultData) >>=
        IO.println("done") >>=
        true
    } else {
        IO.println("abort") >>=
        false
    }
} yield result
```

*Monadic DSL on top of some effect system or Future*

# Asynchronous control flow - Scala industrial usage.

```
for{
    r1 <- cacheService.get(id) match
        case Some(v) => IO.success(v)
        case None => talkToServer(id).map{v =>
            cacheService.update(id,v)
            v
        }
    R2 <- talkToServer(r1.data)
    result <- if (r2.isOk) then {
        writeToFile(r1.data) >>=
        IO.println("done") >>=
        true
    } else {
        IO.println("abort") >>=
        false
    }
} yield result
```

```
async[IO] {
    val r1 = cacheService.getOrUpdate(id,
                                      await(talkToServer(id)))
    val r2 = talkToServer(r1.data)
    if (r2.isOk) then
        writeToFile(r1.data)
        IO.println("done")
        true
    else
        IO.println("abort")
        false
}
```

*"Mainstream" Control-Flow  
over some effect system or Future*

## API (simplified)

```
def async[F[_]](f: T): F[T]
def await[F[_]](f: F[T]): T
```

Well-known `async/await` interface.

- + - generic monad
- ability to use higher-order functions without runtime support for continuations.
- automatic colouring (awaits can-be omitted if can be restored from the context)

## API (simplified)

```
def async[F[_]](f: T):F[T]
def await[F[_]](f: F[T]):T
```

### **transparent inline**

```
def async[F[_]](using am: CpsMonad[F])[T](inline f: am.Context ?=> T):F[T]
    //in real signature here is temporary class, to allow currying of type parameters
```

```
@compileTimeOnly("await should be inside async block")
```

```
def await[G[_],T,F[_]](f: G[T])(using CpsMonadConversion[G,F], CpsMonadContext[F]):
```

## API (less simplified)

### transparent inline

```
def async[F[_]](using am: CpsMonad[F])[T](inline f: am.Context ?=> T): F[T]
```

```
@compileTimeOnly("await should be inside async block")
```

```
def await[G[_], T, F[_]](f: G[T])(using CpsMonadConversion[G, F], CpsMonadContext[F]): T
```

Macro, evaluated during typing

# API (simplified)

**transparent inline**

```
def async[F[_]](using am: CpsMonad[F])[T](inline f: am.Context ?=> T):F[T]
```

```
@compileTimeOnly("await should be inside async block")
```

```
def await[G[_],T,F[_]](f: G[T])(using CpsMonadConversion[G,F], CpsMonadContext[F]): T
```

F[\_]. — Our monad : Future[T], IO[T], ... etc

G[\_]. — Monad which we await

# API (simplified)

```
transparent inline
def async[F[_]](using am: CpsMonad[F])[T](inline f: am.Context ?=> T):F[T]
```

```
@compileTimeOnly("await should be inside async block")
def await[G[_],T,F[_]](f: G[T])(using CpsMonadCoonversion[G,F], CpsMonadContext[F]): T
```

Using implementation of type class,  
available in the current scope

# API (simplified)

```
transparent inline
def async[F[_]](using am: CpsMonad[F])[T](inline f: am.Context ?=> T):F[T]

trait FutureMonad extends CpsMonad[Future] {
  type Context = FutureContex
}

async[Future] {
  // provide API, aviable only inside async
  summon[FutureContext].executionContext.submit(...)
  .....
}
```



Context function

# API (simplified)

**transparent inline**

```
def async[F[_]](using am: CpsMonad[F])[T](inline f: am.Context ?=> T):F[T]
```

```
trait CpsMonad[F[_]] {  
  
  type Context <: CpsMonadContext[F]  
  
  def pure[T](t:T):F[T]  
  def map[A,B](fa:F[A])(f: A=>B):F[B]  
  def flatMap[A,B](fa:F[A])(f: A=>F[B]):F[B]  
  
  def apply[T](op: Context => F[T]): F[T]  
  
}
```

for. try/catch support

```
trait CpsTryMonad[F[_]] extends CpsMonad[F] {  
  
  def error[A](e: Throwable): F[A]  
  def flatMapTry[A,B](fa:F[A])(f: Try[A] => F[B]): F[B]  
}
```

## Transformations: (monadification, based on cps-transform)

rest in continuation passing style



$$C_F[\{a; \underline{\overline{b}}\}]$$

$$\frac{}{F.flatMap(C_F[a])(_ \Rightarrow C_F[b])}$$

# Transformations: control-flow (monadification, based on cps-transform)

// rules for control-flow constructions are straightforward

$$\frac{C_F[\{a; b\}]}{F.flatMap(C_F[a])(\_ \Rightarrow C_F[b])}$$

$$\frac{C_F[val\ a = b; c]}{F.flatMap(C_F[a])(a' \Rightarrow C_F[b_{x/x'}])}$$

$C_F[try\{a\}\{catch\ e \Rightarrow b\}\{finally\ c\}]$

```

F.flatMap(
  F.flatMapTry(C_F[a]){
    case Success(v) => F.pure(v)
    case Failure(e) => C_F[b]
  }
){x => F.map(C_F[c], x)}

```

$$\frac{C_F[t] : t \in Constant \vee t \in Identifier}{F.pure(t)}$$

$$\frac{C_F[if\ a\ then\ b\ else\ c]}{F.flatMap(C_F[a])(a' \Rightarrow if(a')\ then\ C_F[b]\ else\ C_F[c])}$$

$$\frac{C_F[throw\ ex]}{F.error(ex)}$$

---

## Transformations: control-flow, optimisations

Two sequential synchronous fragments are merged into one:

$$\frac{F.flatMap(F.pure(a))(x \Rightarrow F.pure(b(x)))}{F.pure(b(a))}$$

Transformation of each control-flow construction are specialised again sync / async components :

*if a then b else c*

$$C_F[a] \neq F.pure(a) \wedge (C_F[b] \neq F.pure(b) \vee C_F[c] \neq F.pure(c))$$

$$C_F[a] = F.pure(a) \wedge (C_F[b] \neq F.pure(b) \vee C_F[c] \neq F.pure(c))$$

$$C_F[a] = F.pure(a) \wedge C_F[b] = F.pure(b) \wedge C_F[c] = F.pure(c)$$

n(bounds) ~ n(awaits):

Performance == performance of concurrent framework.

# Transformations: higher-order functions

```
class Cache[K, V] {  
  
  def getOrUpdate(k: K, whenAbsent: => V): V  
}
```



Call by name, synonym of  $() \Rightarrow V$

```
async[Process] {  
  val k = retrieveKey(request)  
  cache.getOrUpdate(k, await(fetchValue(k)))  
}
```

$\text{type}(C_F[(x : T)]) = F[T]$ , is not function

$\text{type}(C_F[(x : A \Rightarrow B)]) = A \Rightarrow \text{type}(C_F[B])$

Usual answer: continuation support in runtime environment.

Problem: JVM Runtime (as Js and Native now) have no continuations support

Allow programmer to define async 'shifted' variant of function via typeclass.

# Transformations: higher-order functions

```
class Cache[K,V] {  
    def getOrUpdate(k:K, whenAbsent: =>V): V  
}
```

```
async[Process]{  
    val k = retrieveKey(request)  
    cache.getOrUpdate(k, await(fetchValue(k)))  
}
```

```
summon[AsyncCache[Cache[K,V]]].getOrUpdate[F](cache,monad)(k, ()=>fetchValue(k))  
// where monad = summon[CpsMonad[F]]
```

```
class CacheAsyncShift[K,V] extends  
    AsyncShift[Cache[K,V]]{  
    def getOrUpdate[F[_]](o:Cache[K,V],m:CpsMonad[F])  
        (k:K, whenAbsent: ()=> F[V]):F[V] =  
        ....  
    }
```

# Transformations: higher-order functions

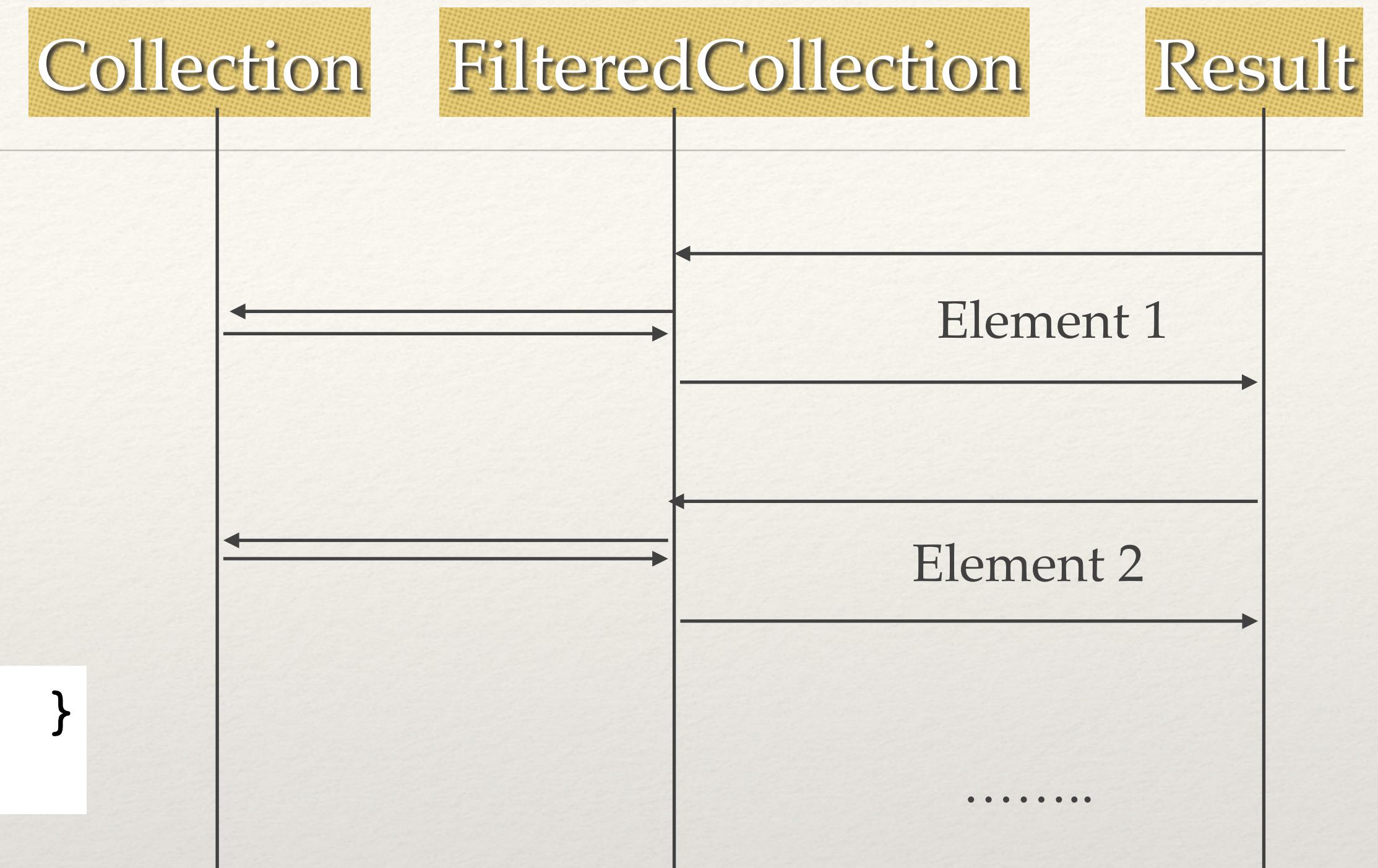
```
class Cache[K,V] {  
    def getOrUpdate(k:K, whenAbsent: =>V): V  
    def getOrUpdateAsync[F[_]](m:CpsMonad[F])(k:K, whenAbsent: ()=>F[V]): F[V]  
}
```

```
async[Process]{  
    val k = retrieveKey(request)  
    cache.getOrUpdate(k, await(fetchValue(k)))  
}
```

```
cache.getOrUpdateAsync[F](monad)(k, ()=>fetchValue(k))  
// where monad = summon[CpsMonad[F]]
```

# substitution classes for chain of higher-order methods

```
class Collection[A] {  
  
    def map(f: A=>B): Collection[B]  
  
    def withFilter(p: A=>Boolean): FilteredCollection[A]  
}  
  
for { url <- urls if await(score(url)) > limit }  
    yield await(fetchData(url))
```



```
urls.withFilter(url => await(score(url))>limit).map(await(fetchData(url)))
```

Developer expect that combination of map and filter will work at the same way, as in synchronous case.

# substitution classes for chain of higher-order methods

```
class Collection[A] {  
    def map(f: A=>B): Collection[B]  
    def withFilter(p: A=>Boolean): FilteredCollection[A]  
}  
  
class AsyncShiftCollectin extends AsyncShift[Collection[A]]:  
  
    def withFilter[F[_]](c:Collection[A],m:CpsMonad[F])  
        (p: A=>F[Boolean]): FilteredCollectionAsyncSubst[F,A]  
  
  
class FilteredCollectionAsyncSubst[F[_],A](m:CpsMonad[F])  
    extends CallChainAsyncShiftSubst[F, FilteredCollection[A], FilteredCollectionAsyncSubst[F,A]]:  
  
    def _finishChain: FilteredCollection[A]  
  
    def withFilter(p: A=> Boolean): this.type  
    def withFilterAsync(p: A=>F[Boolean]): this.type  
    def map[B](f: A => B): this.type  
    def mapAsync[B](f: A=>F[B]): this.type  
  
    // . accumulate all ho-methods in chain, interpret via _finishChain  
    // / categorical interpretation — left Kan extension for FilteredCollection
```

# Transformations:

```
class Collection[A] {  
    def map(f: A=>B): Collection[B]  
    def withFilter(p: A=>Boolean): FilteredCollection[A]  
}  
  
class AsyncShiftCollectin extends AsyncShift[Collection[A]]:  
  
    def withFilter[F[_]](c:Collection[A],m:CpsMonad[F])  
        (p: A=>F[Boolean]): FilteredCollectionAsyncSubst[F,A]  
  
  
class FilteredCollectionAsyncSubst[F[_],A](m:CpsMonad[F])  
    extends CallChainAsyncShiftSubst[F, FilteredCollection[A], FilteredCollectionAsyncSubst[F,A]]:  
  
    def _finishChain: FilteredCollection[A]  
  
    def withFilter(p: A=> Boolean): this.type  
    def withFilterAsync(p: A=>F[Boolean]): this.type  
    def map[B](f: A => B): this.type  
    def mapAsync[B](f: A=>F[B]): this.type  
  
    // . accumulate all ho-methods in chain, interpret via _finishChain  
    // / categorical interpretation — left Kan extension for FilteredCollection
```

# substitution classes for chain of higher-order methods

```
class FilteredCollectionAsyncSubst[F[_], A](m:CpsMonad[F])
  extends CallChainAsyncShiftSubst[F, FilteredCollection[A], FilteredCollecti

class Collection[A] {
  def map(f: A=>B): Collection[B]
  def withFilter(p: A=>Boolean):
}

def _finishChain: FilteredCollection[A]

def withFilter(p: A=> Boolean): this.type
def withFilterAsync(p: A=>F[Boolean]): this.type
def map[B](f: A => B): this.type
def mapAsync[B](f: A=>F[B]): this.type
```

```
urls.withFilter(url => await(score(url))>limit).map(await(fetchData(url)))
```

```
summon[AsynsShift[Collection[A]]]
  .withFilterAsync(url => score(url).map(x => x>limit))
  .mapAsync(fetchData(url))
  ._finishChain()
```

// accumulate all ho-methods in chain, interpret via \_finishChain  
// categorical interpretation — left Kan extension for FilteredCollection

# Transformations: higher-order functions

- ❖ Allow developer to define async variant of sync higher-order methods

- ❖ Externally (in type class)

```
class AsyncShiftCollectin extends AsyncShift[Collection[A]]:

  def withFilter[F[_]](c:Collection[A],m:CpsMonad[F])
    (p: A=>F[Boolean]): FilteredCollectionAsyncSubst[F,A]
```

- ❖ Internally (method with Suffix)

```
class Cache[K,V]:
  def getOrUpdate(k:K, whenAbsent: =>V): V
  def getOrUpdateAsync[F[_]](m:CpsMonad[F])(k:K, whenAbsent: ()=>F[V]): F[V]
```

- ❖ Allow developer to define an async fusion class for evaluation of chain of calls.

```
summon[AsynsShift[Collection[A]]
  .withFilterAsync(url => score(url).map(x => x>limit))
  .mapAsync(fetchData(url))
  ._finishChain()
```

# Automatic colouring.

```
async[IO] {  
    val r1 = cacheService.getOrUpdate(id,  
        await(talkToServer(id)))  
    val r2 = talkToServer(r1.data)  
    if (await(r2).isOk) then  
        await(writeToFile(r1.data))  
        await(IO.println("done"))  
        true  
    else  
        await(IO.println("abort"))  
        false  
}
```



```
async[IO] {  
    val r1 = cacheService.getOrUpdate(id,  
        talkToServer(id))  
    val r2 = talkToServer(r1.data)  
    if (r2.isOk) then  
        writeToFile(r1.data)  
        IO.println("done")  
        true  
    else  
        IO.println("abort")  
        false  
}
```

$x : F[T]) \rightarrow await(x) : T$ . implicit conversion:  $F[T] \Rightarrow T$

discarded value  $F[T] \rightarrow$  discarded value  $await(F[T])$

# Automatic colouring.

```
async[IO] {  
    val r1 = cacheService.getOrUpdate(id,  
        await(talkToServer(id)))  
    val r2 = talkToServer(r1.data)  
    if (await(r2).isOk) then  
        await(writeToFile(r1.data))  
        await(IO.println("done"))  
        true  
    else  
        await(IO.println("abort"))  
        false  
}
```



```
async[IO] {  
    val r1 = cacheService.getOrUpdate(id,  
        talkToServer(id))  
    val r2 = talkToServer(r1.data)  
    if (r2.isOk) then  
        writeToFile(r1.data)  
        IO.println("done")  
        true  
    else  
        IO.println("abort")  
        false  
}
```

// disabled by default, enabled by import

$x : F[T]) \rightarrow await(x) : T$ . implicit conversion:  $F[T] \Rightarrow T$

discarded value  $F[T] \rightarrow$  discarded value  $await(F[T])$

Problem: this can be unsafe for some type of monads.

# Automatic colouring.

$x : F[T] \rightarrow await(x) : T$ . implicit conversion:  $F[T] \Rightarrow T$

discarded value  $F[T] \rightarrow$  discarded value  $await(F[T])$

```
def run(): F[Unit] = async[F] {  
    val connection = openConnection()  
    try  
        while  
            val command = readCommand(connection)  
            val reply = handle(command)  
            if (!reply.isMuted)  
                connection.send(reply.toBytes)  
            !command.isShutdown  
        do ()  
    finally  
        connection.close()  
}
```

This can be unsafe for some type of monads:

```
def run(): F[Unit] = async[F] {  
    val connection = openConnection()  
    try  
        while  
            val command = readCommand(await(connection))  
            val reply = handle(await(command))  
            if (!await(reply).isMuted)  
                await(connection).send(await(reply).toBytes)  
            !await(command).isShutdown  
        do ()  
    finally  
        await(connection).close()  
}
```

$F[\_]$ . Is cached, I.e. two usages refers to the same instance — Safe. (Example: Future)

$F[\_]$ . Is effect, I.e. two usage produce two effect — Unsafe. (Example: IO)

# Automatic colouring.

$x : F[T]) \rightarrow await(x) : T$ . implicit conversion:  $F[T] \Rightarrow T$

discarded value  $F[T] \rightarrow$  discarded value  $await(F[T])$

$F[_]$ . Is effect, I.e. two usage produce two effect — Unsafe. (Example: IO)

$memoize : F[T] \Rightarrow F[F[T]]$

```
def run(): F[Unit] = async[F] {
  val connection = await(memoize(openConnection()))
  try
    while
      val command = await(memoize(readCommand(await(connection))))
      val reply = await(memoize(handle(await(command))))
      if (!await(reply).isMuted)
        await(connection).send(await(reply).toBytes)
      !await(command).isShutdown
    do ()
  finally
    await(connection).close()
}
```

---

## Automatic colouring.

$x : F[T]) \rightarrow await(x) : T$ . implicit conversion:  $F[T] \Rightarrow T$

---

discarded value  $F[T] \rightarrow$  discarded value  $await(F[T])$

$F[\_]$ . Is effect, I.e. two usage produce two effect — Unsafe. (Example: IO)

$memoize : F[T] \Rightarrow F[F[T]]$       Unsafe, when we mix synchronous and asynchronous usage.

Preliminary code analysis two prevent such mixing:

$v : F[T] var, U(v)$  – all usages of  $v$

$u \in U(V)$

Safety condition:

$sync(u) \iff u \sqsubset await(v)$

$async(u) \iff \neg sync(u) \wedge \neg(u \sqsubset (x = v))$

$discarded(u)$

$\neg \exists u_1, u_2 \in U(v) : sync(u_1) \wedge async(u_2)$

$u_2 \sqsubset (v_1 = v_2) \Rightarrow U(v_1) \subset U(v_2)$

---

## Patterns of practical usage:

<https://github.com/rssh/dotty-cps-async>

---

- Migrating from scala2 SIP22 async/await to Scala3 with dotty-cps-async. (Chat-server)
  - Monad: Future
- Utility server for integration of freeswitch with external database.
  - Monad: [X] =>> Resource[IO,X]
  - Stack [ postgres, skunk, http4s ].
- Port of scala library for implementing csp-channels. [scala-gopher]
  - (Non-trivial usage, inline syntax sugar api inside async block)
  - Monad: generic + ReadChannel

### Issues:

- (near 1/2) are compiler issues.
- other: ergonomica and handling specialised cases.

# Loom Impact on JVM.

---

- ❖ Project Loom: <https://wiki.openjdk.org/display/loom/Main>
  - ❖ Original: Continuations for JVM.
  - ❖ Current: Virtual Threads. Continuations interfaces are hidden from the application programmer.
  - ❖ Provides 'await' = 'JavaFuture.get'.
- ❖ Dotty-cps-async - switch for using Loom runtime for suitable monads is in development.
  - ❖ No cps. [future => gibrad mode]
  - ❖ Still need macro for
    - ❖ Awaiting of other monad.
    - ❖ Automatic coloring.
- ❖ Searching for benchmarks. Call me if you want to propose one.



# Embedding Generic Monadic Transformer into Scala.

<https://github.com/rssh/dotty-cps-async>

---

{Questions ? }

*Ruslan Shevchenko, Kyiv, Ukraine.*

*<[ruslan@shevchenko.kiev.ua](mailto:ruslan@shevchenko.kiev.ua)>*

---

*[proofspace.id](http://proofspace.id)*