Liquid Haskell Haskell as a Theorem Prover

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Software bugs are everywhere



Airbus A400M crashed due to a software bug. — May 2015

Software bugs are everywhere



The Heartbleed Bug. Buffer overflow in OpenSSL. 2015

HOW THE HEARTBLEED BUG WORKS:









SERVER, ARE YOU STILL THERE?







Make bugs difficult to express

Using Modern Programming Languages

F#, Ocaml, Erlang, Scala, Haskell

Because of

Strong Types + λ-Calculus

Make bugs difficult to express





VS.





VS.



$\lambda > :m + Data.Text Data.Text.Unsafe$ $<math>\lambda > let pack = "hat"$

λ> :t takeWord16 takeWord16 :: Text -> Int -> Text



$\lambda > :m + Data.Text Data.Text.Unsafe$ $<math>\lambda > let pack = "hat"$

λ> takeWord16 pack True Type Error: Cannot match Bool vs Int



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Valid Values for takeWord16?

takeWord16 :: t:Text -> i:Int -> Text



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Valid Ints	Invalid Ints
0, 1 , len t	len t + 1,

take :: t:Text -> {v:Int | v <= len t} -> Text

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 $\lambda > :m + Data.Text Data.Text.Unsafe$ $<math>\lambda > let pack = "hat"$

λ> take pack 500
Refinement Type Error

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Refinement Type Error

Liquid Haskell

$$len x = 3 \implies v = 500 \implies v <= len x$$

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Checker reports Error

Liquid Haskell: Checks valid arguments, under facts. What are interesting facts?

len "hat" = 3

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(xs ++ ys) ++ zs == xs ++ (ys ++ zs)

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f is == mapReduce f is

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Theorems about Haskell functions

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(Liquid) Haskell as a theorem prover.

Specify theorems as Refinement Types

Prove theorems in Haskell

Use Liquid Haskell to check correctness

data L a = N | C a (L a)

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N ++ ys = ys(C x xs) ++ ys = C x (xs ++ ys)

$$N ++ ys = ys$$

(C x xs) ++ ys = C x (xs ++ ys)

assoc :: xs:L a -> ys:L a -> zs:L a -> ()

$$N ++ ys = ys$$

(C x xs) ++ ys = C x (xs ++ ys)

assoc :: xs:L a -> ys:L a -> zs:L a -> {v:()|(xs ++ ys)++zs==xs++(ys++zs)}

$$N ++ ys = ys$$

(C x xs) ++ ys = C x (xs ++ ys)

assoc :: xs:L a -> ys:L a -> zs:L a -> { (xs ++ ys) ++ zs == xs ++ (ys ++ zs) }

What is the body of assoc?

assoc :: xs:L a -> ys:L a -> zs:L a ->
{ (xs ++ ys) ++ zs == xs ++ (ys ++ zs) }
assoc xs ys zs = ???

(Liquid) Haskell as a theorem prover. What is the body of assoc? A unit Haskell value showing that

left-hand side == right-hand side

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Type Error!

A unit Haskell value showing that left-hand side == right-hand side

assoc :: xs:L a -> ys:L a -> zs:L a ->
{ (xs ++ ys) ++ zs == xs ++ (ys ++ zs) }
assoc xs ys zs
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...
==. xs ++ (ys ++ zs)
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A unit Haskell value showing that left-hand side == right-hand side

assoc :: xs:L a -> ys:L a -> zs:L a -> $\{(xs ++ ys) ++ zs == xs ++ (ys ++ zs)\}$ **assoc** xs ys zs = (xs ++ ys) ++ zs ==. xs ++ (ys ++ zs)*** OED

$$N ++ ys = ys$$

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==. N ++ (ys ++ zs) *** QED

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Thanks!

